



FIDO ECDAA Algorithm

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Editor:

[Rolf Lindemann, Nok Nok Labs, Inc.](#)

Contributors:

[Jan Camenisch, IBM](#)
[Manu Drijvers, IBM](#)
[Alec Edgington, Trustonic](#)
[Anja Lehmann, IBM](#)
[Rainer Urian, Infineon](#)

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Abstract

The FIDO Basic Attestation scheme uses attestation "group" keys shared across a set of authenticators with identical characteristics in order to preserve privacy by avoiding the introduction of global correlation handles. If such an attestation key is extracted from one single authenticator, it is possible to create a "fake" authenticator using the same key and hence indistinguishable from the original authenticators by the relying party. Removing trust for registering new authenticators with the related key would affect the entire set of authenticators sharing the same "group" key. Depending on the number of authenticators, this risk might be unacceptable high.

This is especially relevant when the attestation key is primarily protected against malware attacks as opposed to targeted physical attacks.

An alternative approach to "group" keys is the use of individual keys combined with a Privacy-CA [\[TPMv1-2-Part1\]](#). Translated to FIDO, this approach would require one Privacy-CA interaction for each Uauth key. This means relatively high load and high availability requirements for the Privacy-CA. Additionally the Privacy-CA aggregates sensitive information (i.e. knowing the relying parties the user interacts with). This might make the Privacy-CA an interesting attack target.

Another alternative is the Direct Anonymous Attestation [\[BriCamChe2004-DAA\]](#). Direct Anonymous Attestation is a cryptographic scheme combining privacy with security. It uses the authenticator specific secret once to communicate with a single DAA Issuer and uses the resulting DAA credential in the DAA-Sign protocol with each relying party. The DAA scheme has been adopted by the Trusted Computing Group for TPM v1.2 [\[TPMv1-2-Part1\]](#).

In this document, we specify the use of an improved DAA scheme [\[FIDO-DAA-Security-Proof\]](#) based on elliptic curves and bilinear pairings largely compatible with [\[CheLi2013-ECDAA\]](#) called ECDAA. This scheme provides significantly improved performance compared with the original DAA and basic building blocks for its implementation are part of the TPMv2 specification [\[TPMv2-Part1\]](#).

The improvements over [\[CheLi2013-ECDAA\]](#) mainly consist of security fixes (see [\[ANZ-2013\]](#) and [\[XYZF-2014\]](#)) when splitting the sign operation into two parts.

This specification includes the fixes of the issue regarding (1) the Diffie-Hellman oracle w.r.t. the secret key of the TPM and (2) regarding the potential privacy violations by fraudulent TPMs as proposed in [\[CCDLNU2017-DAA\]](#).

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Table of Contents

- 1. [Notation](#)
 - 1.1 [Conformance](#)
- 2. [Overview](#)

- 2.1 Scope
- 2.2 Architecture Overview
- 3. FIDO ECDAAs Attestation
 - 3.1 Object Encodings
 - 3.1.1 Encoding `BigInteger` values as byte strings (BigIntegerToB)
 - 3.1.2 Encoding `ECPoint` values as byte strings (ECPointToB)
 - 3.1.3 Encoding `ECPoint2` values as byte strings (ECPoint2ToB)
 - 3.2 Global ECDAAs System Parameters
 - 3.3 Issuer Specific ECDAAs Parameters
 - 3.4 ECDAAs-Join
 - 3.4.1 ECDAAs-Join Algorithm
 - 3.4.2 ECDAAs-Join Split between Authenticator and ASM
 - 3.4.3 ECDAAs-Join Split between TPM and ASM
 - 3.5 ECDAAs-Sign
 - 3.5.1 ECDAAs-Sign Algorithm
 - 3.5.2 ECDAAs-Sign Split between Authenticator and ASM
 - 3.5.3 ECDAAs-Sign Split between TPM and ASM
 - 3.6 ECDAAs-Verify Operation
- 4. FIDO ECDAAs Object Formats and Algorithm Details
 - 4.1 Supported Curves for ECDAAs
 - 4.2 ECDAAs Algorithm Names
 - 4.3 `ecdaaSignature` object
- 5. Considerations
 - 5.1 Algorithms and Key Sizes
 - 5.2 Indicating the Authenticator Model
 - 5.3 Revocation
 - 5.4 Pairing Algorithm
 - 5.5 Performance
 - 5.6 Binary Concatenation
 - 5.7 IANA Considerations
- A. References
 - A.1 Normative references
 - A.2 Informative references

1. Notation

Type names, attribute names and element names are written as `code`.

String literals are enclosed in "", e.g. "ED256".

In formulas we use "||" to denote byte wise concatenation operations.

$X = P^x$ denotes scalar multiplication (with scalar x) of a (elliptic) curve point P.

RAND(x) denotes generation of a random number between 0 and x-1.

RAND(G) denotes generation of a random number belonging to Group G.

Specific terminology used in this document is defined in [FIDOGlossary].

The type `BigInteger` denotes an arbitrary length integer value.

The type `ECPoint` denotes an elliptic curve point with its affine coordinates x and y.

The type `ECPoint2` denotes a point on the sextic twist of a BN elliptic curve over $F(q^2)$. The ECPoint2 has two affine coordinates each having two components of type `BigInteger`

1.1 Conformance

As well as sections marked as non-normative, all authoring guidelines, diagrams, examples, and notes in this specification are non-normative. Everything else in this specification is normative.

The key words **must**, **must not**, **required**, **should**, **should not**, **recommended**, **may**, and **optional** in this specification are to be interpreted as described in [RFC2119].

2. Overview

This section is non-normative.

FIDO uses the concept of attestation to provide a cryptographic proof of the **authenticator** [FIDOGlossary] model to the relying party. When the authenticator is registered to the relying party (RP), it generates a new authentication key pair and includes the public key in the attestation message (also known as key registration data object, **KRD**). When using the ECDAAs algorithm, the KRD object is signed using 3.5 ECDAAs-Sign.

For privacy reasons, the authentication key pair is dedicated to one RP (to an application identifier **AppID** [FIDOGlossary] to be more specific). Consequently the attestation method needs to provide the same level of unlinkability. This is the reason why the FIDO ECDAAs Algorithm doesn't use a basename (bsn) often found in other direct anonymous attestation algorithms, e.g. [BriCamChe2004-DAA] or [BFGSW-2011].

The **authenticator** encapsulates all user verification operations and cryptographic functions. An authenticator specific module (**ASM**) [FIDOGlossary] is used to provide a standardized communication interface for authenticators. The authenticator might be implemented in separate hardware or trusted execution environments. The ASM is assumed to run in the normal operating system (e.g. Android, Windows, ...).

2.1 Scope

This document describes the FIDO ECDAAs attestation algorithm in detail.

2.2 Architecture Overview

ECDAAs attestation defines [global system parameters](#) and [ECDAAs issuer specific parameters](#). Both parameter sets need to be installed on the host, in the authenticator and in the FIDO Server. The ECDAAs method consists of two steps:

- [ECDAAs-Join](#) between the authenticator and the **ECDAAs Issuer** to be performed *before* the first FIDO Registration. The [ECDAAs Issuer](#) represents the authenticator vendor as it provides the credentials to attest the authenticator model.
 - $(n, B, sc, yc) = \text{GetNonceFromECDAAsIssuer}()$
 - $(D=Q, c1, s1) = \text{EcdaaSJoin1}(X, Y, B, sc, yc, n)$
 - $(A, B, C, D) = \text{EcdaaSJoin2}(Q, c1, s1)$
 - $\text{EcdaaSJoin2}(A, C) // \text{store cre}=(A, B, C, D)$
- and the pair of [ECDAAs-Sign](#) performed by the authenticator and [ECDAAs-Verify](#) performed by the FIDO Server of the relying party as part of the FIDO Registration.
 - Client: $\text{Attestation} = (\text{signature}, \text{KRD}) = \text{EcdaaSJoin}(\text{AppID})$
 - Server: $\text{success} = \text{EcdaaSVerify}(\text{signature}, \text{KRD}, \text{AppID})$

The technical implementation details of the ECDAAs-Join step are out-of-scope for FIDO. In this document we normatively specify the general algorithm to the extent required for interoperability and we outline examples of some possible implementations for this step.

The ECDAAs-Sign and ECDAAs-Verify steps and the encoding of the related ECDAAs Signature are normatively specified in this document. The generation and encoding of the **KRD** object is defined in other FIDO specifications.

The algorithm and terminology are inspired by [BFGSW-2011](#). The algorithm was modified in order to fix security weaknesses (e.g. as mentioned by [ANZ-2013](#) and [XYZF-2014](#)). Our algorithm proposes an improved task split for the sign operation while still being compatible to TPMv2 (without fixing the TPMv2 weaknesses in such case).

3. FIDO ECDAAs Attestation

This section is normative.

3.1 Object Encodings

We need to convert **BigInteger** and **ECPoint** objects to byte strings using the following encoding functions:

3.1.1 Encoding **BigInteger** values as byte strings (**BigIntegerToB**)

We use the I2OSP algorithm as defined in [RFC3447](#) for converting big numbers to byte arrays. The bytes from the big endian encoded (non-negative) number **n** will be copied right-aligned into the buffer area **b**. The unused bytes will be set to 0. Negative values will not occur due to the construction of the algorithms.

EXAMPLE 1: Converting BigInteger n to byte string b

```
b0 b1 b2 b3 b4 b5 b6 b7
0 0 n0 n1 n2 n3 n4 n5
```

The algorithm implemented in Java looks like this:

EXAMPLE 2: Algorithm for converting BigInteger to byte strings

```
ByteArray BigIntegerToB(
    BigInteger inVal, // IN: number to convert
    int size         // IN: size of the output.
)
{
    ByteArray buffer = new ByteArray(size);
    int oversize = size - inVal.length;
    if (oversize < 0)
        return null;
    for (int i=oversize; i > 0; i--)
        buffer[i] = 0;
    ByteCopy( inVal.bytes, &buffer[oversize], inVal.length);
    return buffer;
}
```

3.1.2 Encoding **ECPoint** values as byte strings (**ECPointToB**)

We use the ANSI X9.62 Point-to-Octet-String [\[ECDSA-ANSI\]](#) conversion using the expanded format, i.e. the format where the compression byte (i.e. 0x04 for expanded) is followed by the encoding of the affine x coordinate, followed by the encoding of the affine y coordinate.

EXAMPLE 3: Converting ECPoint P to byte string

```
(x, y) = ECPointGetAffineCoordinates(P)
len = G1.byteLength
byte string = 0x04 | BigIntegerToB(x,len) | BigIntegerToB(y,len)
```

3.1.3 Encoding **ECPoint2** values as byte strings (**ECPoint2ToB**)

The type **ECPoint2** denotes a point on the sextic twist of a BN elliptic curve over $F(q^2)$, see section [4.1 Supported Curves for ECDAAs](#). Each **ECPoint2** is represented by a pair **(a, b)** of elements of $F(q)$.

The group zero element is always encoded (using the encoding rules as described below) as a an element having all components set to zero (i.e. **cx.a=0, cx.b=0, cy.a=0, cy.b=0**).

We always assume normalized (non-zero) **ECPoint2** values (i.e. **cz = 1**) before encoding them. Non-zero values are encoded using the expanded format (i.e. 0x04 for expanded) followed by the **cx** followed by the **cy** value. This leads to the concatenation of 0x04 followed by the first element (**cx.a**) and second element (**cx.b**) of the pair of **cx** followed by the first element (**cy.a**) and second element (**cy.b**) of the pair of **cy**. All individual numbers are padded to the same length (i.e. the maximum byte length of all relevant 4 numbers).

EXAMPLE 4: Converting ECPoint2 P2 to byte string

```
(cx, cy) = ECPointGetAffineCoordinates(P2)
len = G2.byteLength
byte string = 0x04 | BigIntegerToB(cx.a,len) | BigIntegerToB(cx.b,len)
```

3.2 Global ECDAAs System Parameters

1. Groups G^1 , G^2 and G^T , of sufficiently large prime order p
2. Two generators P^1 and P^2 , such that $G^1 = \langle P^1 \rangle$ and $G^2 = \langle P^2 \rangle$
3. A bilinear pairing $e : G^1 \times G^2 \rightarrow G^T$. We propose the use of "ate" pairing (see [BarNae-2006]). For example source code on this topic, see [BNPairings](#).
4. Hash function H with $H : \{0, 1\}^* \rightarrow Z_p$.
5. Hash function H^{G^1} with $H : \{0, 1\}^* \rightarrow G^1$.
6. $(G^1, P^1, p, H, H^{G^1})$ are installed in all [authenticators](#) implementing FIDO ECDAAs attestation.

Definition of G^1, G^2, G^T , Pairings, hash function H and H^{G^1}

See section [4.1 Supported Curves for ECDAAs](#).

3.3 Issuer Specific ECDAAs Parameters

ECDAAs Issuer Parameters part

1. Randomly generated [ECDAAs Issuer](#) private key $isk = (x, y)$ with $[x, y = RAND(p)]$.
2. [ECDAAs Issuer](#) public key (X, Y) , with $X = P_2^x$ and $Y = P_2^y$.
3. A proof that the [ECDAAs Issuer](#) key was correctly computed
 1. BigInteger $r^x = RAND(p)$
 2. BigInteger $r^y = RAND(p)$
 3. ECPPoint2 $U^x = P_2^{r^x}$
 4. ECPPoint2 $U^y = P_2^{r^y}$
 5. BigInteger $c = H(U^x|U^y|P^2|X|Y)$
 6. BigInteger $s^x = r^x + c \cdot x \pmod{p}$
 7. BigInteger $s^y = r^y + c \cdot y \pmod{p}$
4. $ipk = X, Y, c, s^x, s^y$

Whenever a party uses ipk for the first time, it must first verify that it was correctly generated:

$$H(P_2^{s^x} \cdot X^{-c} | P_2^{s^y} \cdot Y^{-c} | P^2 | X | Y) \stackrel{?}{=} c$$

NOTE

$$P_2^{s^x} \cdot X^{-c} = P_2^{r^x+cx} \cdot P_2^{-cx} = P_2^{r^x} = U^x$$

$$P_2^{s^y} \cdot Y^{-c} = P_2^{r^y+cy} \cdot P_2^{-cy} = P_2^{r^y} = U^y$$

The [ECDAAs Issuer](#) public key ipk **must** be dedicated to a single [authenticator](#) model.

We use the element c of ipk as an identifier for the [ECDAAs Issuer](#) public key (called **ECDAAs Issuer public key identifier**).

3.4 ECDAAs-Join

NOTE

One [ECDAAs-Join](#) operation is required once in the lifetime of an [authenticator](#) prior to the first registration of a credential.

In order to use ECDAAs, the [authenticator](#) must first receive ECDAAs credentials from an [ECDAAs Issuer](#). This is done by the [ECDAAs-Join](#) operation. This operation needs to be performed a single time (before the first credential registration can take place). After the [ECDAAs-Join](#), the [authenticator](#) will use the [ECDAAs-Sign](#) operation as part of each FIDO Registration. The [ECDAAs Issuer](#) is not involved in this step. ECDAAs plays no role in FIDO Authentication / Transaction Confirmation operations.

In order to use ECDAAs, (at least) one [ECDAAs Issuer](#) is needed. The approach specified in this document easily scales to multiple [ECDAAs Issuers](#), e.g. one per [authenticator](#) vendor. FIDO lets the [authenticator](#) vendor choose any [ECDAAs Issuer](#) (similar to his current freedom for selecting any PKI infrastructure/service provider to issuing attestation certificates required for FIDO Basic Attestation).

- All [ECDAAs-Join](#) operations (of the related [authenticators](#)) are performed with one of the [ECDAAs Issuer](#) entities.
- Each [ECDAAs Issuer](#) has a set of public parameters, i.e. [ECDAAs](#) public key material. The related Attestation Trust Anchor is contained in the metadata of each [authenticator](#) model identified by its AAGUID.

There are two different implementation options relevant for the [authenticator](#) vendors (the [authenticator](#) vendor can freely choose them):

1. In-Factory [ECDAAs-Join](#)
2. Remote [ECDAAs-Join](#) and

In the first case, physical proximity is used to locally establish the trust between the [ECDAAs Issuer](#) and the [authenticator](#) (e.g. using a key provisioning

station in a production line). There is no requirement for the ECDAAs Issuer to operate an online web service.

In the second case, some credential is required to remotely establish the trust between the ECDAAs Issuer and the authenticator. As this operation is performed once and only with a single ECDAAs Issuer, privacy is preserved and an authenticator specific credential can and should be used.

Not all ECDAAs authenticators might be able to add their authenticator model IDs (e.g. AAGUID) to the registration assertion (e.g. TPMs). In all cases, the ECDAAs Issuer will be able to derive the exact the authenticator model from either the credential or the physically proximate authenticator. So the ECDAAs Issuer root key must be dedicated to a single authenticator model.

3.4.1 ECDAAs-Join Algorithm

This section is normative.

NOTE

If this join is not in-factory, the value Q must be authenticated by the authenticator. Upon receiving this value, the ECDAAs Issuer must verify that this authenticator did not join before.

1. The authenticator asks the ECDAAs Issuer for the B value of the credential.
2. The ECDAAs Issuer chooses a nonce BigInteger $m = RAND(p)$.
3. The ECDAAs Issuer computes the B value of the credential as $B = HG^1(m)$ and sends $(sc, yc) = HG1_pre(m)$ to the authenticator.
4. The authenticator chooses and stores the ECDAAs private key BigInteger $sk = RAND(p)$
5. The authenticator re-computes $B = H(sc, yc)$
6. The authenticator computes its ECDAAs public key ECPoint $Q = B^{sk}$
7. The authenticator proves knowledge of sk as follows
 1. BigInteger $r^1 = RAND(p)$
 2. ECPoint $U^1 = B^{r^1}$
 3. BigInteger $c^2 = H(U^1|P^1|Q|m)$
 4. BigInteger $n = RAND(p)$
 5. BigInteger $c^1 = H(n|c^2)$
 6. BigInteger $s^1 = r^1 + c^1 \cdot sk$
8. The authenticator sends Q, c^1, s^1, n via the ASM to the ECDAAs Issuer
9. The ECDAAs Issuer verifies that the authenticator is "authentic" and that Q was indeed generated by the authenticator. In the case of an in-factory Join, this might be trivial; in the case of a remote Join this typically requires the use of other cryptographic methods. Since ECDAAs-Join is a one-time operation, unlinkability is not a concern for that.
10. The ECDAAs Issuer verifies that $Q \in G^1$ and verifies $H(n|H(B^{s^1} \cdot Q^{-c^1}|P^1|Q|m)) \stackrel{?}{=} c^1$ (check proof-of-possession of private key).

NOTE

$$B^{s^1} \cdot Q^{-c^1} = B^{r^1+c^1sk} \cdot Q^{-c^1} = B^{r^1+c^1sk} \cdot B^{-c^1sk} = B^{r^1} = U^1$$

11. The ECDAAs Issuer creates credential (A, B, C, D) as follows
 1. ECPoint $A = B^{1/y}$
 2. ECPoint B as computed in the beginning.
 3. ECPoint $C = (A \cdot Q)^x$
 4. ECPoint $D = Q$
12. The ECDAAs Issuer sends A, C to the authenticator. The authenticator still knows B and D
13. The authenticator checks that $A, C \in G^1$ and $A \neq 1^{G^1}$
14. The authenticator checks $e(A, Y) \stackrel{?}{=} e(B, P^2)$

NOTE

$$e(A, Y) = e(B^{1/y}, P^2) = e(B, P^2^{y/y}) = e(B, P^2);$$

15. and the authenticator checks $e(C, P^2) \stackrel{?}{=} e(A \cdot D, X)$

NOTE

$$e(C, P^2) = e((A \cdot Q)^x, P^2); e(A \cdot D, X) = e(A \cdot Q, P^2^x) = e((A \cdot Q)^x, P^2)$$

16. The authenticator stores credential A, B, C, D

3.4.2 ECDAAs-Join Split between Authenticator and ASM

This section is non-normative.

NOTE

If this join is not in-factory, the value Q must be authenticated by the authenticator. Upon receiving this value, the ECDAAs Issuer must verify that this authenticator did not join before.

1. The ASM asks the ECDAAs Issuer for the B value of the credential.
2. The ECDAAs Issuer chooses a nonce BigInteger $m = RAND(p)$
3. The ECDAAs Issuer computes the B value of the credential as $B = HG^1(m)$
4. The ECDAAs Issuer sends $(sc, yc) = HG1_pre(m)$ to the ASM.
5. The ASM forwards (sc, yc) to the authenticator
6. The authenticator chooses and stores the private key BigInteger $sk = RAND(p)$
7. The authenticator re-computes $B = (H(sc), yc)$
8. The authenticator computes its ECDAAs public key ECPoint $Q = B^{sk}$
9. The authenticator proves knowledge of sk as follows
 1. BigInteger $r^1 = RAND(p)$
 2. ECPoint $U^1 = B^{r^1}$
 3. BigInteger $c^2 = H(U^1|P^1|Q|m)$
 4. BigInteger $n = RAND(p)$
 5. BigInteger $c^1 = H(n|c^2)$
 6. BigInteger $s^1 = r^1 + c^1 \cdot sk$
10. The authenticator sends Q, c^1, s^1, n to the ASM, who forwards it to the ECDAAs Issuer.
11. The ECDAAs Issuer verifies that the authenticator is "authentic" and that Q was indeed generated by the authenticator. In the case of an in-factory Join, this might be trivial; in the case of a remote Join this typically requires the use of other cryptographic methods. Since ECDAAs-Join is a one-time operation, unlinkability is not a concern for that.
12. The ECDAAs Issuer verifies that $Q \in G^1$ and verifies $H(n|H(B^{s^1} \cdot Q^{-c^1}|P^1|Q|m)) \stackrel{?}{=} c^1$.
13. The ECDAAs Issuer creates credential (A, B, C, D) as follows
 1. ECPoint $A = B^{1/y}$
 2. ECPoint B as computed in the beginning.
 3. ECPoint $C = (A \cdot Q)^x$
 4. ECPoint $D = Q$
14. The ECDAAs Issuer sends A, C to the ASM. The ASM remembered B and $D = Q$ from an earlier step.
15. The ASM checks that $A, B, C, D \in G^1$ and $A \neq 1^{G^1}$
16. The ASM checks $e(A, Y) \stackrel{?}{=} e(B, P^2)$
17. and the ASM checks that $e(C, P^2) \stackrel{?}{=} e(A \cdot D, X)$
18. The ASM stores A, B, C, D and sends A, C to the authenticator. The authenticator still knows B and D .
19. The authenticator stores B, D and ignores further join requests.

NOTE

These values belong to the ECDAAs secret keys sk . They should persist even in the case of a factory reset.

3.4.3 ECDAAs-Join Split between TPM and ASM

This section is non-normative.

NOTE

The Endorsement key credential (EK-C) and TPM2_ActivateCredentials are used for supporting the remote Join.

This description is based on the principles described in [TPMv2-Part1] section 24 and [Arthur-Challener-2015], page 109 ("Activating a Credential").

1. The ASM asks the ECDAAs Issuer for the B value of the credential.
2. The ECDAAs Issuer chooses a nonce BigInteger $m = RAND(p)$.
3. The ECDAAs Issuer computes the B value of the credential as $B = HG^1(m)$
4. The ECDAAs Issuer sends $(sc, yc) = HG1_pre(m)$ to the ASM.
5. The ASM
 1. instructs the TPM to create a restricted key by calling TPM2_Create, giving the public key template `TPM2_PUBKEY` [TPMv2-Part2] (including the public key P^1 in field `unique`) to the ASM.
 2. re-computes $B = (H(sc), yc)$

3. retrieves TPM Endorsement Key Certificate (EK-C) from the TPM
 4. calls `TPM2_Commit(keyhandle, P1)` where keyhandle is the handle of the restricted key generated before (see above), P1 is set to (B.x,B.y), and s2 and y2 are set to B.x and B.y respectively. This call returns K, E, and ctr; where $K = B^{sk} = Q$, $E = B^{r1}$ is used as U^1 value.
 5. computes BigInteger $c^2 = H(U^1|P^1|Q|m)$
 6. calls `TPM2_Sign(c^2, ctr)`, returning s^1, n .
 7. computes BigInteger $c^1 = H(n|c^2)$
 8. sends EK-C, `TPMT_PUBLIC` (including Q in field `unique`), c^1, s^1, n to the ECDAAs Issuer.
6. The ECDAAs Issuer
1. verifies EK-C and its certificate chain. As a result the ECDAAs Issuer knows the TPM model related to EK-C.
 2. verifies that this EK-C was not used in a (successful) Join before
 3. Verifies that the `objectAttributes` in `TPMT_PUBLIC` [[TPMv2-Part2](#)] matches the following flags: `fixedTPM = 1; fixedParent = 1; sensitiveDataOrigin = 1; encryptedDuplication = 0; restricted = 1; decrypt = 0; sign = 1`.
 4. examines the public key Q, i.e. it verifies that $Q \in G^1$
 5. checks $H(n|H(B^{s^1} \cdot Q^{-c^1}|P^1|Q|m)) \stackrel{?}{=} c^1$
 6. generates the ECDAAs credential (A, B, C, D) as follows
 1. ECPoint $A = B^{1/y}$
 2. ECPoint B as computed in the beginning.
 3. ECPoint $C = (A \cdot Q)^x$
 4. ECPoint $D = Q$
 7. generates a *secret* (derived from a *seed*) and wraps the credential A, B, C, D using that *secret*.
 8. encrypts the *seed* using the public key included in EK-C.
 9. uses *seed* and *name* in KDFa (see [[TPMv2-Part2](#)] section 24.4) to derive HMAC and *symmetric encryption key*. Wrap the *secret* in *symmetric encryption key* and protect it with the *HMAC key*.

NOTE

The parameter *name* in KDFa is derived from `TPMT_PUBLIC`, see [[TPMv2-Part1](#)], section 16.

10. sends the wrapped object including the credential from previous step to the ASM.
7. The ASM instructs the TPM (by calling `TPM2_ActivateCredential`) to
1. decrypt the *seed* using the TPM Endorsement key
 2. compute the *name* (for the ECDAAs attestation key)
 3. use the *seed* in KDFa (with *name*) to derive the *HMAC key* and the *symmetric encryption key*.
 4. use the *symmetric encryption key* to unwrap the *secret*.
8. The ASM
1. unwraps the credential A, B, C, D using the *secret* received from the TPM.
 2. checks that $A, B, C, D \in G^1$ and $A \neq 1G^1$
 3. checks $e(A, Y) \stackrel{?}{=} e(B, P^2)$ and $e(C, P^2) \stackrel{?}{=} e(A \cdot D, X)$
 4. stores A, B, C, D

3.5 ECDAAs-Sign

NOTE

One ECDAAs-Sign operation is required for the client-side environment whenever a new credential is being registered at a relying party.

3.5.1 ECDAAs-Sign Algorithm

This section is normative.

(signature, KRD) = EcdasSign(String AppID)

Parameters

- p: System parameter prime order of group G1 (global constant)
- AppID: FIDO AppID (i.e. https-URL of TrustedFacets object)

Algorithm outline

1. KRD = BuildAndEncodeKRD(); // all traditional Registration tasks are here
2. BigInteger $l = RAND(p)$
3. ECPoint $R = A^l$;
4. ECPoint $S = B^l$;
5. ECPoint $T = C^l$;
6. ECPoint $W = D^l$;

7. BigInteger $r = RAND(p)$
8. ECPPoint $U = S^r$
9. BigInteger $c2 = H(U|S|W|AppID|H(KRD))$
10. BigInteger $n = RAND(p)$
11. $c = H(n | c2)$
12. BigInteger $s = r + c \cdot sk \pmod{p}$
13. signature = (c, s, R, S, T, W, n)
14. return (signature, KRD)

3.5.2 ECDAAsign Split between Authenticator and ASM

This section is non-normative.

NOTE

This split requires both the authenticator and ASM to be honest to achieve anonymity. Only the authenticator must be trusted for unforgeability. The communication between ASM and authenticator must be secure.

Algorithm outline

1. The ASM randomizes the credential
 1. BigInteger $l = RAND(p)$
 2. ECPPoint $R = A^l$;
 3. ECPPoint $S = B^l$;
 4. ECPPoint $T = C^l$;
 5. ECPPoint $W = D^l$;
2. The ASM sends $l, AppID$ to the authenticator
3. The authenticator performs the following tasks
 1. $KRD = BuildAndEncodeKRD()$; // all traditional Registration tasks are here
 2. ECPPoint $S' = B^l$
 3. ECPPoint $W' = D^l$
 4. BigInteger $r = RAND(p)$
 5. ECPPoint $U = S^r$
 6. BigInteger $c2 = H(U|S'|W'|AppID|H(KRD))$
 7. BigInteger $n = RAND(p)$
 8. $c = H(n | c2)$
 9. BigInteger $s = r + c \cdot sk \pmod{p}$
 10. Send c, s, KRD, n to the ASM
4. The ASM sets signature = (c, s, R, S, T, W, n) and outputs (signature, KRD)

3.5.3 ECDAAsign Split between TPM and ASM

This section is non-normative.

NOTE

This algorithm is for the special case of a TPMv2 as authenticator. This case requires both the TPM and ASM to be honest for anonymity. Only the TPM must be trusted for unforgeability (see [CCDLNU2017-DAA]).

Algorithm outline

1. The ASM randomizes the credential
 1. BigInteger $l = RAND(p)$
 2. ECPPoint $R = A^l$;
 3. ECPPoint $S = B^l$;
 4. ECPPoint $T = C^l$;
 5. ECPPoint $W = D^l$;
2. The ASM calls TPM2_Commit() with $P1$ set to S and $s2, y2$ empty buffers. The ASM receives the result values $K, L, E = S^r = U$ and ctr. K and L are empty since $s2, y2$ are empty buffers.
3. The ASM calls TPM2_Create to generate the new authentication key pair. The related private key might need to be protected with appropriate access control mechanisms, e.g. see section 8 of [UAFAuthnrCommands].
4. The ASM calls TPM2_Certify() on the newly created key with ctr from the TPM2_Commit and $E = U, S, W, AppID$ as qualifying data. The ASM receives signature value s and related nonce n and attestation block KRD (i.e. TPMS_ATTEST structure in this case).
5. BigInteger $c2 = H(E|S|W|AppID|H(KRD))$, using KRD as returned by the previous step.
6. The ASM computes: $c = H(n | c2)$

7. The ASM sets signature = (c, s, R, S, T, W, n) and outputs (signature, KRD)

3.6 ECDAAs-Verify Operation

This section is normative.

NOTE

One ECDAAs-Verify operation is required for the FIDO Server as part of each FIDO Registration.

boolean EcdasVerify(signature, AppID, KRD, ModelName)

Parameters

- p: System parameter prime order of group G^1 (global constant)
- P^2 : System parameter generator of group G^2 (global constant)
- signature: (c, s, R, S, T, W, n)
- AppID: FIDO AppID
- KRD: Attestation Data object as defined in other specifications.
- ModelName: the claimed FIDO authenticator model (i.e. either AAID or AAGUID)

Algorithm outline

1. Based on the claimed ModelName, look up X, Y from trusted source
2. Check that $R, S, T, W \in G^1$, $R \neq 1^{G^1}$, and $S \neq 1^{G^1}$.
3. $H(n|H(S^s \cdot W^{-c}|S|W|AppID|H(KRD))) \stackrel{?}{=} c$; fail if not equal

NOTE

$$B = A^y = P_1^{ly}$$

$$D = Q^{ly} = P_1^{skly} = B^{sk}$$

$$S = B^l \text{ and } W = D^l$$

$$U = S^r$$

$$\begin{aligned} S^s \cdot W^{-c} &= S^{r+csk} \cdot W^{-c} = U \cdot S^{csk} \cdot W^{-c} \\ &= U \cdot B^{lcsk} \cdot D^{-lc} = U \cdot B^{lcsk} \cdot B^{-lcsk} = U \end{aligned}$$

4. $e(R, Y) \stackrel{?}{=} e(S, P^2)$; fail if not equal

NOTE

$$e(R, Y) = e(A^l, P_2^y); e(S, P^2) = e(B^l, P^2) = e(A^{ly}, P^2)$$

5. $e(T, P^2) \stackrel{?}{=} e(R \cdot W, X)$; fail if not equal

NOTE

$$e(T, P^2) = e(C^l, P^2) = e(A^{xl} \cdot Q^{xlyl^j}, P^2); e(A^l \cdot D^l, X) = e(A^l \cdot Q^{lyl^j}, P_2^x)$$

6. for (all sk' on RogueList) do if $W \stackrel{?}{=} S^{sk'}$ fail;
7. // perform all other processing steps for new credential registration

NOTE

In the case of a TPMv2, i.e. KRD is a TPMS_ATTEST object. In this case the verifier must check whether the TPMS_ATTEST object starts with TPM_GENERATED magic number and whether its field objectAttributes contains the flag fixedTPM=1 (indicating that the key was generated by the TPM).

8. return true;

4. FIDO ECDAAs Object Formats and Algorithm Details

This section is normative.

4.1 Supported Curves for ECDAAs

Definition of G1

G1 is an elliptic curve group $E : y^2 = x^3 + ax + b$ over $F(q)$ with $a = 0$.

Definition of G2

G2 is the p-torsion subgroup of $E'(Fq^2)$ where E' is a sextic twist of E. With $E' : y'^2 = x'^3 + b'$.

An element of $F(q^2)$ is represented by a pair (a,b) where $a + bX$ is an element of $F(q)[X]/\langle X^2 + 1 \rangle$. We use angle brackets $\langle Y \rangle$ to signify the ideal generated by the enclosed value.

NOTE

In the literature the pair (a,b) is sometimes also written as a complex number $a + b * i$.

Definition of GT

GT is an order-p subgroup of Fq^{12} .

Pairings

We propose the use of Ate pairings as they are efficient (more efficient than Tate pairings) on Barreto-Naehrig curves [DevScoDah2007].

Supported BN curves

We use pairing-friendly Barreto-Naehrig [BarNae-2006] [ISO15946-5] elliptic curves. The curves `TPM_ECC_BN_P256` and `TPM_ECC_BN_P638` curves are defined in [TPMv2-Part4].

BN curves have a Modulus $q = 36 \cdot u^4 + 36 \cdot u^3 + 24 \cdot u^2 + 6 \cdot u + 1$ [ISO15946-5] and a related order of the group $p = 36 \cdot u^4 + 36 \cdot u^3 + 18 \cdot u^2 + 6 \cdot u + 1$ [ISO15946-5].

- `TPM_ECC_BN_P256` is a curve of form $E(F(q))$, where q is the field modulus [TPMv2-Part4] [BarNae-2006]. This curve is identical to the P256 curve defined in [ISO15946-5] section C.3.5.
 - The values have been generated using $u = -7\ 530\ 851\ 732\ 716\ 300\ 289$.
 - Modulus $q = 115\ 792\ 089\ 237\ 314\ 936\ 872\ 688\ 561\ 244\ 471\ 742\ 058\ 375\ 878\ 355\ 761\ 205\ 198\ 700\ 409\ 522\ 629\ 664\ 518\ 163$
 - Group order $p = 115\ 792\ 089\ 237\ 314\ 936\ 872\ 688\ 561\ 244\ 471\ 742\ 058\ 035\ 595\ 988\ 840\ 268\ 584\ 488\ 757\ 999\ 429\ 535\ 617\ 037$
 - p and q have length of 256 bit each.
 - $b = 3$
 - $P^1_{256} = (x=1, y=2)$
 - $b' = (a=3, b=3)$
 - $P^2_{256} = (x,y)$, with
 - $P^2_{256}.x = (a=114\ 909\ 019\ 869\ 825\ 495\ 805\ 094\ 438\ 766\ 505\ 779\ 201\ 460\ 871\ 441\ 403\ 689\ 227\ 802\ 685\ 522\ 624\ 680\ 861\ 435, b=35\ 574\ 363\ 727\ 580\ 634\ 541\ 930\ 638\ 464\ 681\ 913\ 209\ 705\ 880\ 605\ 623\ 913\ 174\ 726\ 536\ 241\ 706\ 071\ 648\ 811)$
 - $P^2_{256}.y = (a=65\ 076\ 021\ 719\ 150\ 302\ 283\ 757\ 931\ 701\ 622\ 350\ 436\ 355\ 986\ 716\ 727\ 896\ 397\ 520\ 706\ 509\ 932\ 529\ 649\ 684, b=113\ 380\ 538\ 053\ 789\ 372\ 416\ 298\ 017\ 450\ 764\ 517\ 685\ 681\ 349\ 483\ 061\ 506\ 360\ 354\ 665\ 554\ 452\ 649\ 749\ 368)$
- `TPM_ECC_BN_P638` [TPMv2-Part4] uses
 - The values have been generated using $u = 365\ 375\ 408\ 992\ 443\ 362\ 629\ 982\ 744\ 420\ 548\ 242\ 302\ 862\ 098\ 433$
 - Modulus $q = 641\ 593\ 209\ 463\ 000\ 238\ 284\ 923\ 228\ 689\ 168\ 801\ 117\ 629\ 789\ 043\ 238\ 356\ 871\ 360\ 716\ 989\ 515\ 584\ 497\ 239\ 494\ 051\ 781\ 991\ 794\ 253\ 619\ 096\ 481\ 315\ 470\ 262\ 367\ 432\ 019\ 698\ 642\ 631\ 650\ 152\ 075\ 067\ 922\ 231\ 951\ 354\ 925\ 301\ 839\ 708\ 740\ 457\ 083\ 469\ 793\ 717\ 125\ 223$
 - The related order of the group is $p = 641\ 593\ 209\ 463\ 000\ 238\ 284\ 923\ 228\ 689\ 168\ 801\ 117\ 629\ 789\ 043\ 238\ 356\ 871\ 360\ 716\ 989\ 515\ 584\ 497\ 239\ 494\ 051\ 781\ 991\ 794\ 253\ 619\ 096\ 481\ 315\ 470\ 262\ 367\ 432\ 019\ 698\ 642\ 631\ 650\ 152\ 075\ 067\ 922\ 231\ 951\ 354\ 925\ 301\ 839\ 708\ 740\ 457\ 083\ 469\ 793\ 717\ 125\ 222, y=16)$
 - p and q have length of 638 bit each.
 - $b = 257$
 - $P^1_{638} = (x=641\ 593\ 209\ 463\ 000\ 238\ 284\ 923\ 228\ 689\ 168\ 801\ 117\ 629\ 789\ 043\ 238\ 356\ 871\ 360\ 716\ 989\ 515\ 584\ 497\ 239\ 494\ 051\ 781\ 991\ 794\ 253\ 619\ 096\ 481\ 315\ 470\ 262\ 367\ 432\ 019\ 698\ 642\ 631\ 650\ 152\ 075\ 067\ 922\ 231\ 951\ 354\ 925\ 301\ 839\ 708\ 740\ 457\ 083\ 469\ 793\ 717\ 125\ 222, y=16)$
 - $b' = (a=771, b=1542)$
 - $P^2_{638} = (x, y)$, with
 - $P^2_{638}.x = (a=192\ 492\ 098\ 325\ 059\ 629\ 927\ 844\ 609\ 092\ 536\ 807\ 849\ 769\ 208\ 589\ 403\ 233\ 289\ 748\ 474\ 758\ 010\ 838\ 876\ 457\ 636\ 072\ 173\ 883\ 771\ 602\ 089\ 605\ 233\ 264\ 992\ 910\ 618\ 494\ 201\ 909\ 695\ 576\ 234\ 119\ 413\ 319\ 303\ 931\ 909\ 848\ 663\ 554\ 062\ 144\ 113\ 485\ 982\ 076\ 866\ 968\ 711\ 247, b=166\ 614\ 418\ 891\ 499\ 184\ 781\ 285\ 132\ 766\ 747\ 495\ 170\ 152\ 701\ 259\ 472\ 324\ 679\ 873\ 541\ 478\ 330\ 301\ 406\ 623\ 174\ 002\ 502\ 345\ 930\ 325\ 474\ 988\ 134\ 317\ 071\ 869\ 554\ 535\ 111\ 092\ 924\ 719\ 466\ 650\ 228\ 182\ 095\ 841\ 246\ 668\ 361\ 451\ 788\ 368\ 418\ 036\ 777\ 197\ 454\ 618\ 413\ 255)$
 - $P^2_{638}.y = (a=622\ 964\ 952\ 935\ 200\ 827\ 531\ 506\ 751\ 874\ 167\ 806\ 262\ 407\ 152\ 244\ 280\ 323\ 674\ 626\ 687\ 789\ 202\ 660\ 794\ 092\ 633\ 841\ 098\ 984\ 322\ 671\ 973\ 226\ 667\ 873\ 503\ 889\ 270\ 602\ 870\ 064\ 426\ 165\ 592\ 237\ 410\ 681\ 318\ 519\ 893\ 784\ 898\ 821\ 343\ 051\ 339\ 820\ 566\ 224\ 981\ 344\ 169\ 470, b=514\ 285\ 963\ 827\ 225\ 043\ 076\ 463\ 721\ 426\ 569\ 583\ 576\ 029\ 220\ 880\ 138\ 564\ 906\ 219\ 230\ 942\ 887\ 639\ 456\ 599\ 654\ 554\ 743\ 732\ 087\ 558\ 187\ 149\ 207\ 036\ 952\ 474\ 092\ 411\ 405\ 629\ 612\ 957\ 921\ 369\ 286\ 372\ 038\ 525\ 830\ 610\ 755\ 207\ 588\ 843\ 864\ 366\ 759\ 521\ 090\ 861\ 911\ 494)$
- `ECC_BN_DSD_P256` [DevScoDah2007] section 3 uses
 - The values have been generated using $u = 6\ 917\ 529\ 027\ 641\ 089\ 837$
 - Modulus $q = 82434016654300679721217353503190038836571781811386228921167322412819029493183$
 - The related order of the group is $p = 82434016654300679721217353503190038836284668564296686430114510052556401373769$
 - p and q have length of 256 bit each.
 - $b = 3$

- $P^1_{DSD_P256} = (1, 2)$
- $b' = (a=3, b=6)$
- $P^2_{DSD_P256} = (x, y)$, with
 - $P^2_{DSD_P256.x} = (a=73\ 481\ 346\ 555\ 305\ 118\ 071\ 940\ 904\ 527\ 347\ 990\ 526\ 214\ 212\ 698\ 180\ 576\ 973\ 201\ 374\ 397\ 013\ 567\ 073\ 039, b=28\ 955\ 468\ 426\ 222\ 256\ 383\ 171\ 634\ 927\ 293\ 329\ 392\ 145\ 263\ 879\ 318\ 611\ 908\ 127\ 165\ 887\ 947\ 997\ 417\ 463)$
 - $P^2_{DSD_P256.y} = (a=3\ 632\ 491\ 054\ 685\ 712\ 358\ 616\ 318\ 558\ 909\ 408\ 435\ 559\ 591\ 759\ 282\ 597\ 787\ 781\ 393\ 534\ 962\ 445\ 630\ 353, b=60\ 960\ 585\ 579\ 560\ 783\ 681\ 258\ 978\ 162\ 498\ 088\ 639\ 544\ 584\ 959\ 644\ 221\ 094\ 447\ 372\ 720\ 880\ 177\ 666\ 763)$
- **ECC_BN_ISOP512** [ISO15946-5] section C.3.7 uses
 - The values have been generated using $u=138\ 919\ 694\ 570\ 470\ 098\ 040\ 331\ 481\ 282\ 401\ 523\ 727$
 - Modulus $q = 13\ 407\ 807\ 929\ 942\ 597\ 099\ 574\ 024\ 998\ 205\ 830\ 437\ 246\ 153\ 344\ 875\ 111\ 580\ 494\ 527\ 427\ 714\ 590\ 099\ 881\ 795\ 845\ 981\ 157\ 516\ 604\ 994\ 291\ 639\ 750\ 834\ 285\ 779\ 043\ 186\ 149\ 750\ 164\ 319\ 950\ 153\ 126\ 044\ 364\ 566\ 323$
 - The related order of the group is $p = 13\ 407\ 807\ 929\ 942\ 597\ 099\ 574\ 024\ 998\ 205\ 830\ 437\ 246\ 153\ 344\ 875\ 111\ 580\ 494\ 527\ 427\ 714\ 590\ 099\ 881\ 680\ 053\ 891\ 920\ 200\ 409\ 570\ 720\ 654\ 742\ 146\ 445\ 677\ 939\ 306\ 408\ 461\ 754\ 626\ 647\ 833\ 262\ 056\ 300\ 743\ 149$
 - p and q have length of 512 bit each.
 - $b = 3$
 - $P^1_{ISO_P512} = (x=1, y=2)$
 - $b' = (a=3, b=3)$
 - $P^2_{ISO_P512} = (x, y)$, with
 - $P^2_{ISO_P512.x} = (a=3\ 094\ 648\ 157\ 539\ 090\ 131\ 026\ 477\ 120\ 117\ 259\ 896\ 222\ 920\ 557\ 994\ 037\ 039\ 545\ 437\ 079\ 729\ 804\ 516\ 315\ 481\ 514\ 566\ 156\ 984\ 245\ 473\ 190\ 248\ 967\ 907\ 724\ 153\ 072\ 490\ 467\ 902\ 779\ 495\ 072\ 074\ 156\ 718\ 085\ 785\ 269, b=3\ 776\ 690\ 234\ 788\ 102\ 103\ 015\ 760\ 376\ 468\ 067\ 863\ 580\ 475\ 949\ 014\ 286\ 077\ 855\ 600\ 384\ 033\ 870\ 546\ 339\ 773\ 119\ 295\ 555\ 161\ 718\ 985\ 244\ 561\ 452\ 474\ 412\ 673\ 836\ 012\ 873\ 126\ 926\ 524\ 076\ 966\ 265\ 127\ 900\ 471\ 529)$
 - $P^2_{ISO_P512.y} = (a=7\ 593\ 872\ 605\ 334\ 070\ 150\ 001\ 723\ 245\ 210\ 278\ 735\ 800\ 573\ 263\ 881\ 411\ 015\ 285\ 406\ 372\ 548\ 542\ 328\ 752\ 430\ 917\ 597\ 485\ 450\ 360\ 707\ 892\ 769\ 159\ 214\ 115\ 916\ 255\ 816\ 324\ 924\ 295\ 339\ 525\ 686\ 777\ 569\ 132\ 644\ 242, b=9\ 131\ 995\ 053\ 349\ 122\ 285\ 871\ 305\ 684\ 665\ 648\ 028\ 094\ 505\ 015\ 281\ 268\ 488\ 257\ 987\ 110\ 193\ 875\ 868\ 585\ 868\ 792\ 041\ 571\ 666\ 587\ 093\ 146\ 239\ 570\ 057\ 934\ 816\ 183\ 220\ 992\ 460\ 187\ 617\ 700\ 670\ 514\ 736\ 173\ 834\ 408)$

NOTE

Spaces are used inside numbers to improve readability.

Hash Algorithm H

Depending on the curve, we use $H(x) = \text{SHA256}(x) \bmod p$ or $H(x) = \text{SHA512}(x) \bmod p$ as hash algorithm $H: \{0, 1\}^* \rightarrow \mathbb{Z}_p$.

The argument of the hash function must always be converted to a byte string using the appropriate encoding function specific in section 3.1 Object Encodings, e.g. according to section 3.1.3 Encoding ECPoint2 values as byte strings (ECPoint2ToB) in the case of `ecPoint2` points.

NOTE

We don't use [IEEE P1363.3](#) section 6.1.1 IHF1-SHA with security parameter t (e.g. $t=128$ or 256) as it is more complex and not supported by TPMv2.

Hash Algorithm H^{G^1}

Definition of H^{G^1} (taken from [CheLi2013-ECDA]):

$H^{G^1}: \{0, 1\}^* \rightarrow G^1$, where G^1 is an elliptic curve group $E: y^2 = x^3 + b$ over $\text{GF}(q)$ with cofactor = 1. Given a message $m \in \{0, 1\}^*$,

H^{G^1} can be computed as follows:

ECPoint $p = \text{HG1}(\text{String } m)$

1. Set $i = 0$ be a 32-bit unsigned integer.
2. Compute $x = H(\text{BigNumberToB}(i, 4) \parallel m)$
3. Compute $z = x^3 + b \bmod q$
4. Compute $y = \text{sqrt}(z) \bmod q$. If y does not exist, set $i = i + 1$, repeat step 2 if $i < 232$, otherwise, report failure.
5. Set $y = \min(y, q - y)$.
6. return `ECPoint(x, y)`

(String sc , BigNumber yc) = $\text{HG1_pre}(\text{String } m)$

1. Set $i = 0$ be a 32-bit unsigned integer.
2. Compute $x = H(\text{BigNumberToB}(i, 4) \parallel m)$
3. Compute $z = x^3 + b \bmod q$
4. Compute $y = \text{sqrt}(z) \bmod q$. If y does not exist, set $i = i + 1$, repeat step 2 if $i < 232$, otherwise, report failure.
5. Set $y = \min(y, q - y)$.
6. Set sc to `BigNumberToB}(i, 4) \parallel m`.
7. Set yc to y .
8. return (sc, yc)

The ASM on the FIDO User device platform can help the authenticator compute $\text{HG1}(m)$, yet the authenticator verifies the computation as follows: Given m , the ASM runs the above algorithm. For a successful execution, let $sc = (\text{istr } m)$ and yc be the y value in the last step. The ASM sends sc and

yc to the authenticator. The authenticator computes $HG1(m) = (H(sc), yc)$.

Given the value sc, the original message m can be recomputed by skipping the first 4 bytes.

4.2 ECDAAs Algorithm Names

We define the following JWS-style algorithm names (see [RFC7515]):

ED256

[TPM_ECC_BN_P256](#) curve, using SHA256 as hash algorithm H.

ED256-2

[ECC_BN_DSD_P256](#) curve, using SHA256 as hash algorithm H.

ED512

[ECC_BN_ISOP512](#) curve, using SHA512 as hash algorithm H.

ED638

[TPM_ECC_BN_P638](#) curve, using SHA512 as hash algorithm H.

4.3 ecdaaSignature object

The fields c and s both have length N. The fields R, S, T, W have equal length ($2*N+1$ each).

In the case of BN_P256 curve (with key length N=32 bytes), the fields R, S, T, W have length $2*32+1=65$ bytes. The fields c and s have length N=32 each.

The ecdaaSignature object is a binary object generated as the concatenation of the binary fields in the order described below (total length of 356 bytes for 256bit curves):

Value	Length (in Bytes)	Description
UINT8[] ECDAAsignature_c	N	The c value, $c = H(n c2)$ as returned by EcdaaSign encoded as byte string according to BigIntegerToB. Where <ul style="list-style-type: none"> $c2 = H(U S W KR AppID)$ $U = S^r$, with $r = RAND(p)$ computed by the signer. KR is the the entire to-be-signed object (e.g. TAG_UAFV1_KR in the case of FIDO UAF). $S = B^l$, with $l = RAND(p)$ computed by the signer and $B = A^y$ computed in the ECDAAs-Join
UINT8[] ECDAAsignature_s	N	The s value, $s = r + c * sk \pmod p$, as returned by EcdaaSign encoded as byte string according to BigIntegerToB. Where <ul style="list-style-type: none"> $r = RAND(p)$, computed by the signer at FIDO registration (see 3.5.2 ECDAAs-Sign Split between Authenticator and ASM) p is the group order of G1 sk: is the authenticator's attestation secret key, see above
UINT8[] ECDAAsignature_n	N	The Nonce value n, as returned by EcdaaSign encoded as byte string according to BigIntegerToB.
UINT8[] ECDAAsignature_R	$2*N+1$	$R = A^l$; computed by the ASM or the authenticator at FIDO registration; encoded as byte string according to ECPointToB. Where <ul style="list-style-type: none"> $l = RAND(p)$, i.e. random number $0 \leq l < p$. Computed by the ASM or the authenticator at FIDO registration. And where $R = A^l$ denotes the scalar multiplication (of scalar l) of a curve point A. Where A has been provided by the ECDAAs Issuer as part of ECDAAs-Join: $A = B^{1/y}$, see 3.4.1 ECDAAs-Join Algorithm. Where p is a system value, injected into the authenticator and y is part of the ECDAAs Issuer private key $isk=(x,y)$.
UINT8[] ECDAAsignature_S	$2*N+1$	$S = B^l$; computed by the ASM or the authenticator at FIDO registration encoded as byte string according to ECPointToB. Where B has been provided by the ECDAAs Issuer on Join: $B = HG1(m) = (H(sc), yc)$, see 3.4.1 ECDAAs-Join Algorithm .
UINT8[] ECDAAsignature_T	$2*N+1$	$T = C^l$; computed by the ASM or the authenticator at FIDO registration encoded as byte string according to ECPointToB. Where <ul style="list-style-type: none"> $C = (A \cdot Q)^x$, provided by the ECDAAs Issuer on Join x is a components of the ECDAAs Issuer private key, $isk=(x,y)$. Q is the authenticator public key
UINT8[] ECDAAsignature_W	$2*N+1$	$W = D^l$; computed by the ASM or the authenticator at FIDO registration encoded as byte string according to ECPointToB. Where $D = Q$ is computed by the ECDAAs Issuer at Join (see 3.4.1 ECDAAs-Join Algorithm).

Value	Length (in Bytes)	Description
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5. Considerations

This section is non-normative.

A detailed security analysis of this algorithm can be found in [\[FIDO-DAA-Security-Proof\]](#).

5.1 Algorithms and Key Sizes

The proposed algorithms and key sizes are chosen such that compatibility to TPMv2 is possible.

5.2 Indicating the Authenticator Model

Some authenticators (e.g. TPMv2) do not have the ability to include their model (i.e. vendor ID and model name) in attested messages (i.e. the to-be-signed part of the registration assertion). The TPM's endorsement key certificate typically contains that information directly or at least it allows the model to be derived from the endorsement key certificate.

In FIDO, the relying party expects the ability to cryptographically verify the authenticator model.

We require the ECDAAs Issuers public key (ipk=(X,Y,c,sx,sy)) to be dedicated to one single authenticator model (e.g. as identified by AAID or AAGUID).

5.3 Revocation

If the private ECDAAs attestation key *sk* of an authenticator has been leaked, it can be revoked by adding its value to a RogueList.

The ECDAAs-Verifier (i.e. FIDO Server) check for such revocations. See section [3.6 ECDAAs-Verify Operation](#).

The ECDAAs Issuer is expected to check revocation by other means:

1. if ECDAAs-Join is done in-factory, it is assumed that produced devices are known to be uncompromised (at time of production).
2. if a remote ECDAAs-Join is performed, the (remote)ECDAAs Issuer already must use a different method to remotely authenticate the authenticator (e.g. using some endorsement key). We expect the ECDAAs Issuer to perform a revocation check based on that information. This is even more flexible as it does not require access to the authenticator ECDAAs private key *sk*.

5.4 Pairing Algorithm

The pairing algorithm [e](#) needs to be used by the ASM as part of the Join process and by the verifier (i.e. FIDO relying party) as part of the verification (i.e. FIDO registration) process.

The result of such a pairing operation is only compared to the result of another pairing operation computed by the same entity. As a consequence, it doesn't matter whether the ASM and the verifier use the exact same pairings or not (as long as they both use valid pairings).

5.5 Performance

For performance reasons the calculation of $Sig2=(R, S, T, W)$ may be performed by the ASM running on the FIDO user device (as opposed to inside the authenticator). See section [3.5.2 ECDAAs-Sign Split between Authenticator and ASM](#).

The cryptographic computations to be performed inside the authenticator are limited to G1. The ECDAAs Issuer has to perform two G2 point multiplications for computing the public key. The Verifier (i.e. FIDO relying party) has to perform G1 operations and two pairing operations.

5.6 Binary Concatentation

We use a simple byte-wise concatenation function for the different parameters, i.e. $H(a,b) = H(a || b)$.

This approach is as secure as the underlying hash algorithm since the authenticator controls the length of the (fixed-length) values (e.g. U, S, W). The AppID is provided externally and has unverified structure and length. However, it is only followed by a fixed length entry - the (system defined) hash of KRD. As a consequence, no parts of the AppID would ever be confused with the fixed length value.

5.7 IANA Considerations

This specification registers the algorithm names "ED256", "ED512", and "ED638" defined in section [4. FIDO ECDAAs Object Formats and Algorithm Details](#) with the IANA JSON Web Algorithms registry as defined in section "Cryptographic Algorithms for Digital Signatures and MACs" in [\[RFC7518\]](#).

Algorithm Name	"ED256"
Algorithm Description	FIDO ECDAAs algorithm based on TPM_ECC_BN_P256 [TPMv2-Part4] curve using SHA256 hash algorithm.
Algorithm Usage Location(s)	"alg", i.e. used with JWS.
JOSE Implementation Requirements	Optional
Change Controller	FIDO Alliance, Contact Us
Specification Documents	Sections 3. FIDO ECDAAs Attestation and 4. FIDO ECDAAs Object Formats and Algorithm Details of [FIDOecdaasAlgorithm] .
Algorithm Analysis Document(s)	[FIDO-DAA-Security-Proof]

Algorithm Name	"ED512"
Algorithm Description	ECDAAs algorithm based on ECC_BN_ISOP512 [ISO15946-5] curve using SHA512 algorithm.
Algorithm Usage Location(s)	"alg", i.e. used with JWS.
JOSE Implementation Requirements	Optional
Change Controller	FIDO Alliance, Contact Us

Specification Documents	Sections 3. FIDO ECDA A Attestation and 4. FIDO ECDA A Object Formats and Algorithm Details of [FIDOEcdaaAlgorithm] .
Algorithm Analysis Document(s)	[FIDO-DAA-Security-Proof]
Algorithm Name	"ED638"
Algorithm Description	ECDA A algorithm based on TPM_ECC_BN_P638 [TPMv2-Part4] curve using SHA512 algorithm.
Algorithm Usage Location(s)	"alg", i.e. used with JWS.
JOSE Implementation Requirements	Optional
Change Controller	FIDO Alliance, Contact Us
Specification Documents	Sections 3. FIDO ECDA A Attestation and 4. FIDO ECDA A Object Formats and Algorithm Details of [FIDOEcdaaAlgorithm] .
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